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Quantum Correlations as Necessary Precondition for Secure Communication

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- Interface Physics Computer Science in Quantum Communication
 - Physics provides correlations with a promise
 - Computer Science uses correlations within complex communication task
- Classical and Quantum Correlations
 - If Physics is to add something, then we need correlations with quantum features
- 'Entanglement' as necessary conditions for quantum communication
- Exploitation of conditions
 - entanglement witnesses
 - application to 6-state, 4-state and 2-state protocol (QKD)
- Conclusions



Bennett Brassard Protocol





Quantum Communication and Correlations

Phase I: Physical Set-Up

Generation of correlations between Alice and Bob

→ possibly containing hidden correlations with Eve

Physics: correlated data with a promise.

Which type of correlations are

useful for Quantum Communication?

(Classical) Computer Science:

Solve Communication Problem with classically correlated data ...

Phase II: Classical Communication Protocol

Advantage distillation (e.g. announcement of bases in BB84 protocol) Error Correction (→ Alice and Bob share the same key) Privacy Amplification (→ generates secret key shared by Alice and Bob) Note: classical communication for QKD can be improved:

e.g. in QKD with weak light pulses [Acín, Gisin, and Scarani, Phys. Rev. A 69, 012309 (2004)] or two-way communication [Lo, Gottesman, guant-ph/0105121]





Key extraction from correlated classical data

Lower bound on secrecy capacity C_S: (rate of secret communication between Alice and Bob) - Csiszar, Körner, IEEE, IT 24, 339 (1978).

$$C_{S} > \max \{I_{AB} - I_{AE}, I_{AB} - I_{BE}\}$$

Derived for classical three-party correlations Eve: quantum system! I. Devetak, A. Winter, quant-ph/0307053.

<u>Upper Bounds on secrecy capacity C_S:</u>

- U. M. Maurer, IEEE Trans. Inf. Theo. 39, 1733 (1993); -U. Maurer and S. Wolf, IEEE T. I. T. **45**, 499 (1999).

$$Cs \le I(A; B \downarrow E)$$

• Intrinsic Information: I(A;B↓E)

 $I(A,B\downarrow E) = \min_{E\to \underline{E}} I(A;B|\underline{E}) \text{ with } I(A;B|E) = H(A,E) + H(B,E) - H(A,B,E) - H(E)$

<u>Quantum</u>

$$P_{F}(A,B,E) = P(a,b) \operatorname{Tr} (\rho_{E}(a,b) F_{E})$$

$$I(A;B\downarrow E) = \inf_{F} I_{F}(A,B|E)$$

'Information' Bob can gain about Alice's data by looking at his own data, whatever Eve told him about Alice's data.



Intercept/Resend attack



→P(a,b,e)=p(a,e) p(b|e) (Markov Chain)
→ Intrinsic information vanishes, no secret communication possible!

Example:

BB84 with

- •Poissonian photon number distribution
- •losses in the quantum channel
- •symmetric error rate in signals
- → implementing **specific** intercept/resend



[M.Curty, N.L, in preparation]

(for vanishing error rate:

[Jahma, Dusek, NL, Phys. Rev. A 62, 022306 (2000)]





Are these correlations useful?

Assumptions:

trusted ideal source of ideal BB84 protocol \checkmark

trusted ideal detector of ideal BB84 protocol

 \times

Probability Distribution P(A,B)

	0	1	+	-
0	0.07987	0.04516	0.00913	0.11591
1	0.04508	0.07986	0.11593	0.00901
+	0.11599	0.00909	0.08001	0.04507
-	0.00897	0.11593	0.04505	0.07985

Error Rate: 36%



Entanglement behind the scene

How to generate correlated classical data:







Necessary condition for secure communication

Knowledge available to Alice and Bob:

measurement POVM {A_i}_i, {B_j}_j (may contain imperfections!)
observed joint probability distribution P(A,B)
[red. density matrix ρ_A (P&M schemes)]

Theorem (Entanglement Based and P&M):

• If P(A,B) together with $\{A_i\}_i$, $\{B_j\}_j$ [and ρ_A for P&M schemes] allows interpretation as separable state then I(A;B \downarrow E) = 0, and therefore C_S = 0.

M. Curty., M. Lewenstein and N. L, quant-ph/0307151.

Theorem: (converse)

• $I(A;B\downarrow E) > 0$ iff P(A,B) together with $\{A_i\}_i$, $\{B_j\}_j$ cannot be interpreted as coming from a separable state.

-A. Acín and N. Gisin, quant-ph/0310054.

NOTE: does not guarantee a secret key ...

Observation of quantum correlation excludes intercept/resend attack!

Approach allows for realistic implementations! -detection inefficiency goes into {B_i}_i

-full mode description of sender and receiver



Entanglement verification

Problem structure:

- \bullet Unknown density matrix ρ_{AB}
- constraints via observed correlations (data) P(A,B) [for P&M schemes: fixed ρ_A]
- Question: any separable ρ_{AB} compatible with constraints?

Specific experiment and data:

search for entanglement proof (sufficient, not necessary)

- rule out separability e.g. via Bell inequality
- violation of local uncertainty relations [Hofmann, Takeuchi, PRA 68 032103 (2003)]
- numerical optimisation via entanglement witnesses [Eisert, Hyllus, Gühne, Curty, quant-ph/040713

Specific experiment:

- general efficient numerical method for any possible data?
- find analytic complete necessary and sufficient condition for any possible data
 - \rightarrow approach in following part for simple qubit protocols



Entanglement Witnesses

Entangled States:

 ρ_{AB} is entangled iff $\rho_{AB} \neq \Sigma_i p_i |a_i\rangle\langle a_i|_A \otimes |b_i\rangle\langle b_i|_B$

Entanglement Witnesses (EW):

• ρ_{AB} is entangled iff \exists W hermitian such that:

$$\begin{split} &Tr\{W{\cdot}\rho_{AB}\} < 0 \\ &Tr\{W{\cdot}\sigma_{AB}\} \ge 0 \ \forall \ \sigma_{AB} \text{ separable} \end{split}$$

-M. Horodecki, P. Horodecki and R. Horodecki, Phys. Lett. A **223**, 1 (1996). -M.B. Terhal, Phys. Lett. A **271**, 319 (2000).

Optimal EW (OEW):

-M. Lewenstein, B. Kraus, J.I. Cirac and P. Horodecki, PRA 62, 052310 (2000).







Local Measurement of Entanglement Witnesses

Decomposition of Witnesses in Local Measurements:

Any bipartite hermitian operator W can be decomposed as a *pseudo-mixture*:

$$W = \Sigma_{ij} c_{ij} A_i \otimes B_j \qquad \text{with} \qquad c_{ij} \in \mathfrak{R}, \ \Sigma_{ij} c_{ij} = 1$$

where $A_i \otimes B_i$ forms a POVM operator basis.

 \rightarrow {A_i}_i, {B_j}_j describe measurements (positive, add up to identiy) <u>Evaluation</u>:

Then

Tr{W·
$$\rho_{AB}$$
} = $\Sigma_{ij} c_{ij} Tr{A_i \otimes B_j \rho_{AB}} = \Sigma_{ij} c_{ij} P(a_i, b_j)$

- -O. Gühne, P. Hyllus, D. Bruss, A. Ekert, M. Lewenstein, C. Macchiavello and A. Sanpera, PRA 66, 062305 (2002).
- -O. Gühne, P. Hyllus, D. Bruss, A. Ekert, M. Lewenstein, C. Macchiavello and A. Sanpera, J. Mod. Opt. 50 (6-7), 1079 (2003).

⁻A. Sanpera, R. Tarrach and G. Vidal, PRA 58, 826 (1997).



Necessary condition based on entanglement witnesses

Theorem:

• Given a set of local operations with POVM elements $A_i \bigotimes B_j$ together with the probability distribution of their ocurrence, P(A,B), then the correlations P(A,B) cannot lead to a secret key via public communication unless one can prove the presence of entenglement in the (effectively) distributed state via an entanglement witnesses $W = \sum_{ij} c_{ij} A_i \bigotimes B_j$ with c_{ij} real such that $Tr\{W\sigma_{AB}\} \ge 0$ for all separable states σ_{AB} and $\sum_{ij} c_{ij} P(i,j) < 0$.

-M. Curty, M. Lewenstein and N. L., Phys. Rev. Lett. 92, 217903 (2004).

Important point:

entanglement witness criterion is necessary and sufficient even for restricted knowledge about the shared quantum state!

Idea:

states with verifiable entanglement form a convex set →restricted class of witnesses can testify the verifiable entanglement compatible with sep. verifyable entangled • ρ_{AB}



Use three mutually unbiased bases: e.g. X,Y,Z direction in Bloch sphere

- Bruß, Phys. Rev. 81, 3018 (1998);
- Bechmann-Pasquinucci et al, PRA 59, 4238 (1999) .

Simplified thought experiment: use two-qubit state:



6-State QKD protocol

$$W_6 = \sum_{ij} c_{ij} \sigma_i \otimes \sigma_j$$

with
$$\mathbf{i}, \mathbf{j} = \{0, \mathbf{x}, \mathbf{z}, \mathbf{y}\}$$
, and $\sigma_0 = 1$.

- Include all Optimal DEW: $W = |\psi_e\rangle\langle\psi_e|^{T_B}$
- All entangled states can be detected.

Searching for quantum correlations:

- parametrize $|\psi_e\rangle$
- evaluate locally Tr[$\rho |\psi_e \rangle \langle \psi_e |^{T_B}$]
- search for negative expectation values



Use two mutually unbiased bases: e.g. X,Z direction in Bloch sphere

-C.H. Bennett and G. Brassard, Proc. IEEE Int. Conf. On Computers, System and Signal Processing, 175 (1984).

Observation:

4-State QKD protocol

$$W_4^{EBS} = \sum_{ij} c_{ij} \sigma_i \otimes \sigma_j$$

with
$$i, j = \{0, x, z\}$$
, and $\sigma_0 = 1$.

 \rightarrow restricted class of witnesses

$$W \in W_4^{EBS}$$
 iff $W = W^T = W^{T_B}$

- Alice and Bob cannot evaluate Optimal DEW.
- Not all entangled states can be detected.



4-State QKD protocol

Optimal W_4^{EBS} (OEW₄) verifiable entangled with sep. ρ_{AB} $W \in OEW_4^{EBS}$

Observation:

Given $W \in W_4^{EBS}$ necessary to detect entanglement in state ρ_{AB} is that the operator

 $\Omega = \rho_{AB} + \rho_{AB}^{T} + \rho_{AB}^{T} + \rho_{AB}^{T}$ is a non-positive operator.

<u>Theorem</u>: The EW that are optimal within the four-state protocol are given by

$$OEW_4^{EBS} = \frac{1}{2}(Q+Q^{T_B})$$

with $Q = |\psi_e\rangle\langle\psi_e|$ such that $Q = Q^T$

-M. Curty., M. Lewenstein and N. L., quant-ph/0307151.

• OEW_4^{EBS} provides necessary and sufficient conditions for detection of quantum correlations in P(A,B).

• For P&M schemes we find $OEW_4^{P&M} = OEW_4^{EBS}$



Quantum Correlations? (II)

Assumptions: (BB84 setup) trusted ideal source

trusted ideal detector

Probability Distribution P(A,B)

A∖B	0	1	+	-			
0	0.07987	0.04516	0.00913	0.11591			
1	0.04508	0.07986	0.11593	0.00901			
+	0.11599	0.00909	0.08001	0.04507			
-	0.00897	0.11593	0.04505	0.07985			
Error Rate: 36 %							



values are marked)

Witness Class:

 $OEW_4^{EBS} = \frac{1}{2} \left(|\psi_e\rangle \langle \psi_e| + |\psi_e\rangle \langle \psi_e|^{T_B} \right)$

 $|\psi_{e}\rangle = \cos(X)|00\rangle + \sin(X)(\cos(Y)|01\rangle + \sin(Y)(\cos(Z)|10\rangle + \sin(Z)|11\rangle))$





Use two non-orthogonal states, e.g., $|\phi_0\rangle$ and $|\phi_1\rangle$

-C.H. Bennett, Phys. Rev. Lett. 68, 3121 (1992).

<u>Theorem:</u>

The family

$$W_2 = |0\rangle\langle 0|\otimes A + |1\rangle\langle 1|\otimes B + x C(\theta)$$

with $A = A^T$, $B = B^T$, $A \ge 0$, $B \ge 0$, rank(A) = rank(B) = 2, $\theta \in [0, 2\pi)$, and

 $\mathbf{x} = \min_{|\phi\rangle} (\langle \phi | \mathbf{A} | \phi \rangle \langle \phi | \mathbf{B} | \phi \rangle)^{1/2}$

is sufficient to detect all entangled states that are detectable in the 2-state protocol.

-M. Curty., O. Gühne, M. Lewenstein and N. L, (in preparation).

2-State EW:

$$W_{2} = \sum_{i} c_{i} \sigma_{0} \otimes \sigma_{i} + \sum_{j} c_{j} \sigma_{z} \otimes \sigma_{j} + \sum_{k} c_{k} \sigma_{k} \otimes \sigma_{0}$$

with $i, j = \{x, z\}, k = \{0, x, z, y\}, and \sigma_0 = 1$.

 \rightarrow restricted class of witnesses





iterface Physics – Computer Science: Classical Correlated Data with a Promise

ecessary condition for secure QKD is the proof of presence of quantum correlations

uantum correlations: for entanglement based <u>and</u> prepare&measure schemes.

For experiments: show the presence of such entanglement

- •no need to enter details of classical communication protocols
- •prevents oversights in preliminary analyses
- •one properly constructed entanglement proof (e.g entanglement witness) suffices

For theory:

- show in which situation quantum correlations are **sufficient** to generate secret key
- develop figure of merit (secrecy capacity) to measure secrecy potential of correlations.
- develop proper entanglement proofs for realistic experiments (for given measurements)
- develop **compact description for restricted class of entanglement witnesses** (allows effective search of quantum correlations)
- include **detection inefficiencies** into the witness construction